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# Theory of Hypermatroids (Applied Combinatorial Theory and Algorithms)

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Theory of Hypermatroids

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Summary

Hypermatroid is a generalized notion of matroid and network flow. However, it is not a mere generalization but it gives us a deep insight into matroids and network flows and also provides us with new significant problems which are overlooked. This is an abstract of the series of the author's papers [1] ~ [7].

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If  $E$  is a finite set, we let  $\mathfrak{S}_0(E)$  denote the linear space of real-valued *modular functions* on  $2^E$  which corresponds closely to the vector space  $R^E$  except on  $\phi$ . Let  $\mu^\phi$  denote the *constant function* in  $\mathfrak{S}_0(E)$  such that  $\mu^\phi(X) = 1$  for any  $X \subseteq E$  and let  $\mu^i$  denote the *unit function* such that  $\mu^i(X) = 1$  if  $i \in X$  and 0 otherwise. For any  $\xi \in \mathfrak{S}_0(E)$ , we define

$$\text{car}^\pm \xi = \{i \mid \xi(\{i\}) \gtrless \xi(\phi)\}.$$

A *hedron*  $\mathfrak{B}$  in  $\mathfrak{S}_0(E)$  is a compact non-empty subset of  $\mathfrak{S}_0(E)$  satisfying the following

Exchange axiom for bases [1]. If  $\xi, \eta \in \mathfrak{B}$  and  $\xi \neq \eta$ , then  $\xi(\phi) = \eta(\phi)$  and  $\exists i \in \text{car}^-(\xi - \eta)$ ,  $\exists j \in \text{car}^+(\xi - \eta)$ ,  $\exists \hat{c} > 0$ ;

$$(0 \leq) \forall c \leq \hat{c}: \xi + c(\mu^i - \mu^j), \eta + c(\mu^j - \mu^i) \in \mathfrak{B}.$$

We say the modular function  $\xi \in \mathfrak{B}$  is a *base* of  $\mathfrak{B}$ . Since the Exchange axiom for bases has a self-dual form, it is obvious that if  $\mathfrak{B}$  is a hedron in  $\mathfrak{S}_0(E)$  then  $-\mathfrak{B} \equiv \{-\xi \mid \xi \in \mathfrak{B}\}$  is also a hedron in  $\mathfrak{S}_0(E)$ , called the *inverse* of  $\mathfrak{B}$ .

The *hyperspace* (or *hedron space*)  $\mathfrak{P}(E)$  is the linear space of all hedra in  $\mathfrak{S}_0(E)$  such that the *sum* is defined by

$$\mathfrak{A} + \mathfrak{B} \equiv \{\xi + \eta \mid \xi \in \mathfrak{A}, \eta \in \mathfrak{B}\}$$

and *scalar multiple* is defined by

$$c\mathfrak{B} \equiv \{c\eta \mid \eta \in \mathfrak{B}\} \quad (c \in \mathbb{R}).$$

Obviously,  $\mathfrak{S}_0(E)$  is a subspace of  $\mathfrak{P}(E)$ . Thus a hyperspace can be regarded as a natural generalization of a vector space.

Let  $\mathfrak{S}_\pm(E)$  denote the convex cone of real-valued *super[sub]-modular functions* on  $2^E$ . The *deficiency [rank] function*  $\sigma_\pm$  of a hedron  $\mathfrak{B}$  is defined by

$$\sigma_\pm(X) \equiv \min_{\max} \{ \xi(X) \mid \xi \in \mathfrak{B} \}.$$

Then we have the following

Theorem 1. The deficiency [rank] function  $\sigma_\pm$  of a hedron  $\mathfrak{B}$  in  $\mathfrak{P}(E)$  is a super[sub]modular function, that is  $\sigma_\pm \in \mathfrak{S}_\pm(E)$ .

The remarkable property of Theorem 1 is that its converse is also true.

Theorem 2. For any super[sub]modular function  $\sigma_\pm \in \mathfrak{S}_\pm(E)$ , the convex polyhedron  $\mathfrak{B}_\pm$  in  $\mathfrak{S}_0(E)$  defined by

$$\mathfrak{B}_\pm \equiv \{ \xi \mid \xi \geq \sigma_\pm, \xi(E) = \sigma_\pm(E), \xi(\phi) = \sigma_\pm(\phi) \}$$

is a hedron in  $\mathfrak{P}(E)$ .

We have then established a one-to-one correspondence between the hedra in  $\mathfrak{P}(E)$  and the super[sub]modular functions in  $\mathfrak{S}_\pm(E)$ .

The concept of *hypermattroids* may be defined in several different but equivalent ways. That is, a hypermatroid  $\mathfrak{M}$  is a

pair  $(E, \mathfrak{B})$ ,  $(E, \sigma_+)$ ,  $(E, \sigma_-)$ , etc., where  $E$  is a ground set,  $\mathfrak{B}$  is a hedron in  $\mathfrak{P}(E)$ , and  $\sigma_{\pm} (\in \mathfrak{S}_{\pm}(E))$  is the deficiency [rank] function of  $\mathfrak{B}$ .

The  $\ast$ -dual of  $\sigma_{\pm} \in \mathfrak{S}_{\pm}(E)$  is defined by

$$\sigma_{\pm}^*(X) \equiv -\sigma_{\pm}(\bar{X}) + \sigma_{\pm}(E) + \sigma_{\pm}(\phi).$$

Theorem 3. The  $\ast$ -dual of the deficiency [rank] function of a hedron  $\mathfrak{B}$  in  $\mathfrak{P}(E)$  is the rank [deficiency] function of  $\mathfrak{B}$ , that is  $\sigma_{\pm}^* = \sigma_{\mp}$ .

Hereafter, for notational simplicity let  $\rho$  be a submodular function on  $2^E$ . Define a *least upper vector*  $\hat{\rho}$  and a *greatest lower vector*  $\check{\rho}$  of  $\rho$  by

$$\hat{\rho}(x) \equiv \rho(\{x\}) - \rho(\phi) \quad (\forall x \in E),$$

$$\check{\rho}(x) \equiv \rho(E) - \rho(E - \{x\}) \quad (\forall x \in E),$$

respectively. And define a *least upper modular function*  $\hat{\rho}$  and a *greatest lower modular function*  $\check{\rho}$  of  $\rho$  by

$$\hat{\rho}(X) \equiv \sum \hat{\rho}(x) \quad (x \in X),$$

$$\check{\rho}(X) \equiv \sum \check{\rho}(x) \quad (x \in X),$$

respectively. We say  $\rho^{\circ} \equiv \hat{\rho} - \check{\rho}$  ( $\in \mathfrak{S}_0(E)$ ) is the *oscillation* of  $\rho$ . The  $\flat$ -dual of  $\rho$  is defined by  $\rho^{\flat} \equiv \rho^{\circ} - \rho$  and the  $\sharp$ -dual of  $\rho$  is defined by  $\rho^{\sharp} \equiv (\rho^{\flat})^*$ . Then we call the hypermatroid  $\mathfrak{M}^{\sharp} =$

$(E, \rho^h)$  defined by a rank function the *dual* of  $\mathfrak{M} = (E, \rho)$ . The hedron  $\mathfrak{B}^h$  of  $\mathfrak{M}^h$  is given by

$$\mathfrak{B}^h \equiv \{\rho^o - \xi \mid \xi \in \mathfrak{B}\},$$

and the deficiency function of  $\mathfrak{M}^h$  is given by  $\rho^b$ . Obviously we have

$$(\mathfrak{M}^h)^h = \mathfrak{M}.$$

Note that the duality of hypermatroids is slightly different from the duality of matroids.

Theorem 4. Any submodular function  $\rho$  can be decomposed into three parts as follows:

$$\rho = \rho(\phi)\mu^\phi + \check{\rho} + \tilde{\rho}. \quad (1)$$

We call  $\tilde{\rho}$  in (1) the *proper* submodular function of  $\rho$ . Let  $\tilde{\mathfrak{S}}_-(E)$  denote the set of all proper submodular functions in  $\mathfrak{S}_-(E)$ .

A hypermatroid  $\mathfrak{M} = (E, \rho)$  is called *integral* if  $\rho$  is integer-valued. A *polymatroid* [8]  $\mathfrak{M} = (E, \rho)$  is a hypermatroid satisfying  $\rho(\phi) = 0$  and  $\hat{\rho}(x) \geq 0$  ( $\forall x \in E$ ). A *matroid* is an integral polymatroid satisfying  $\hat{\rho}(x) \leq 1$  ( $\forall x \in E$ ).

A *quasimatroid* [7] is an integral polymatroid such that its rank function is the direct sum of proper submodular functions

satisfying

$$\left. \begin{aligned} \rho(X \cup \{y\}) + \rho(X \cup \{z\}) - \rho(X) - \rho(X \cup \{y, z\}) &= 0, 1 \\ (y, z \notin X, y \neq z), \end{aligned} \right\} (2)$$

unit functions and the constant function. Obviously a matroid is a quasimatroid.

The following theorem solves the open question by Edmonds [8].

Extreme rays theorem [2]. The extreme rays of  $\tilde{\mathcal{E}}_-(E)$  are the proper submodular functions which satisfy the above condition (2) and have minimal sets of intervals  $[X, X \cup \{y, z\}]$  such that the left hand side of (2) is equal to 1.

Let  $N = (V, A; c)$  be the capacitated network, where  $V$  is a vertex set,  $A$  is a directed arc set and  $c$  is a capacity vector in  $R_+^A$ . Define the *cut function*  $\gamma: 2^V \rightarrow R_+$  by

$$\gamma(X) \equiv c(X, \bar{X}) \equiv \sum c(a) \quad (\partial^+ a \in X, \partial^- a \in \bar{X}, a \in A).$$

For any *flow*  $f \in R_+^A$  satisfying  $f \leq c$ , define the *boundary function*  $\partial f: 2^V \rightarrow R$  by

$$\partial f(X) \equiv \sum \{f(\delta^+ v) - f(\delta^- v)\} \quad (v \in X).$$

Then  $\gamma$  is submodular and we have  $\partial f(\emptyset) = \gamma(\emptyset) = 0$ ,  $\partial f(V) = \gamma(V) = 0$  and  $\partial f \leq \gamma$ . Therefore,  $\mathfrak{R} = (V, \gamma)$  is a hypermatroid defined by

a rank function and every boundary function  $\partial f$  for a flow  $f$  is a base of  $\mathfrak{N}$  [5],[6]. Thus we have known that a capacitated network  $N = (V, A; c)$  is a typical example of a hypermatroid.

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